Combinatorial link Floer homology and transverse knots

Dylan Thurston

Joint with/work of Sucharit Sarkar
Ciprian Manolescu
Peter Ozsváth
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Lenhard Ng

math.GT/{0607691,0610559,0611841,0703446}
http://www.math.columbia.edu/~dpt/speaking

June 10, 2007, Princeton, NJ

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The invariant called knot Heegaard-Floer
Determines the genus–and more.
To distinguish transverse knots
(and it turns out there are lots!)
HFK opens up a new door.

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Outline

- **Introduction**

Computing *HFK*

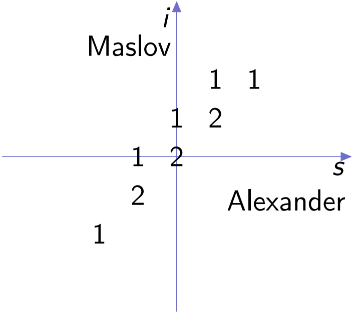
Variants

Grid moves

Transverse knots

What is Heegaard-Floer homology?

$\dim(\widehat{HFK}_i(K;s))$:



Characteristics of $\widehat{HFK}$:

- Bigraded;
- Euler characteristic is Conway-Alexander polynomial;
- Max grading is knot genus; (Ozsváth-Szabó 2001)
- Determines knot fibration; (Ghiggini, Ni 2006)
- Defined via pseudo-holomorphic curves.

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...and so the world's simplest algorithm for knot genus!

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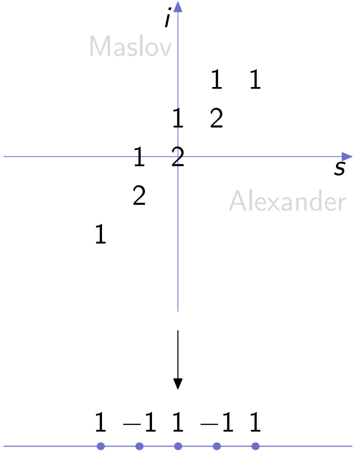

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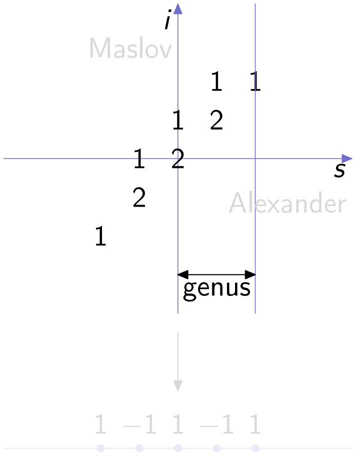

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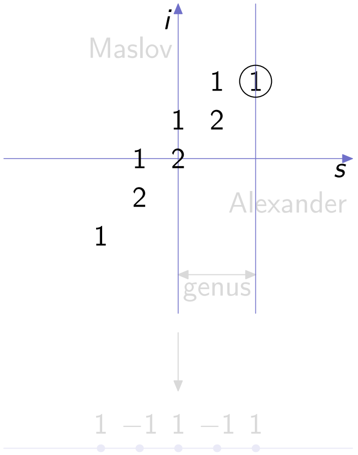

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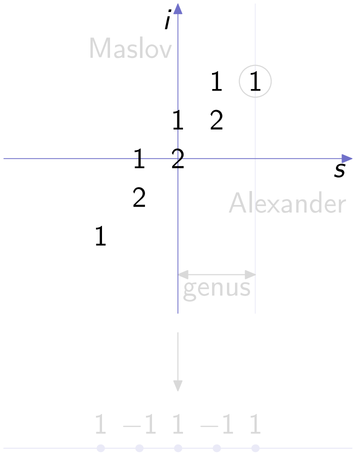

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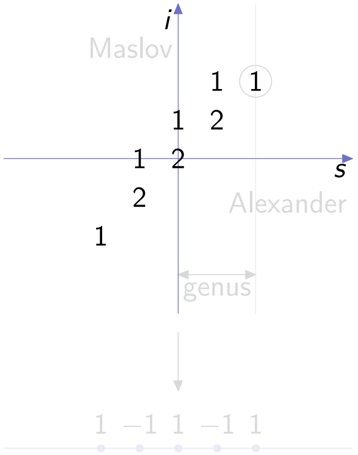

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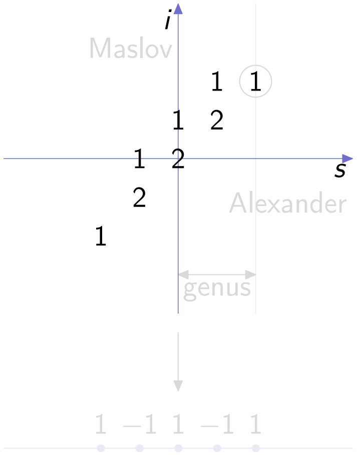

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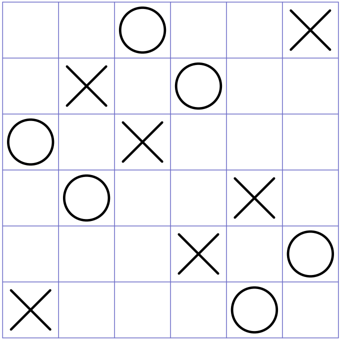
Setting: Grid diagrams



Grid diagram: square diagram with one X and one O per row and column.

Turn it into a knot: connect

X to O in each column;

O to X in each row.

Cross vertical strands over horizontal.

Grid diagrams exist: take any diagram, rotate crossings so vertical crosses over horizontal.

The knot is unchanged under *cyclic rotations*:

Move top segment to bottom.

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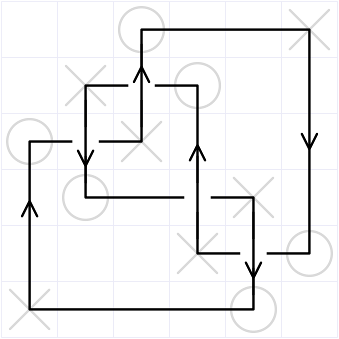

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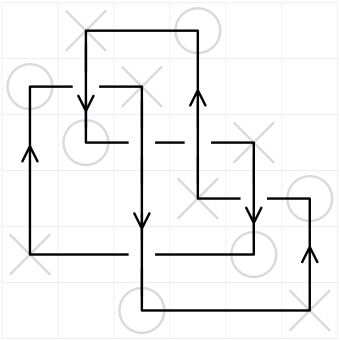

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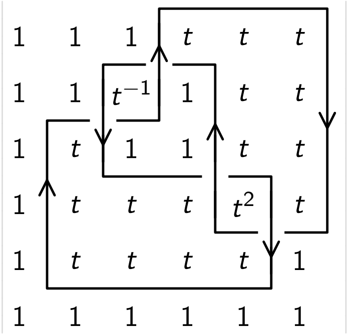
Computing the Alexander polynomial

We categorify the following formula:



$$\begin{vmatrix} 1 & 1 & 1 & t & t & t \\ 1 & 1 & t^{-1} & 1 & t & t \\ 1 & t & 1 & 1 & t & t \\ 1 & t & t & t & t^2 & t \\ 1 & t & t & t & t & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 \end{vmatrix} = \pm t^*(1-t)^{n-1}\Delta(K;t)$$

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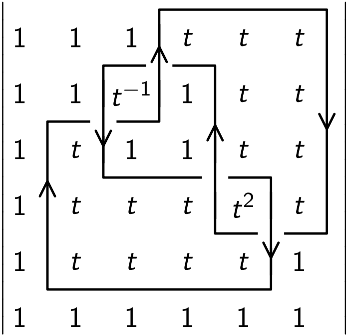

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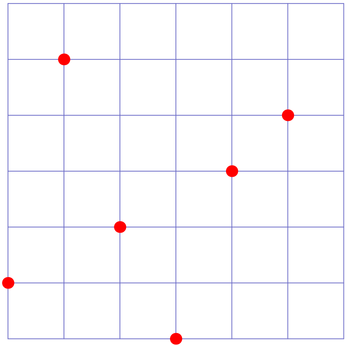
Computing *HFK*: Chain complex $\widetilde{CK}$



Define a chain complex $\widetilde{CK}$ over $\mathbb{Z}/2$.

- Generated by matchings between horizontal and vertical gridcircles (as counted in det for Alexander).
- Boundary ∂ switches corners on *empty rectangles*:

Sum over all ways to switch SW-NE corners of an empty rectangle to NW-SE corners. (*Empty* means: no X's, O's, or other points in generator.)

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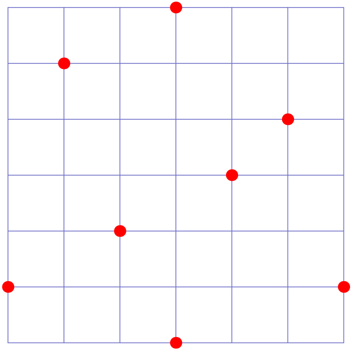

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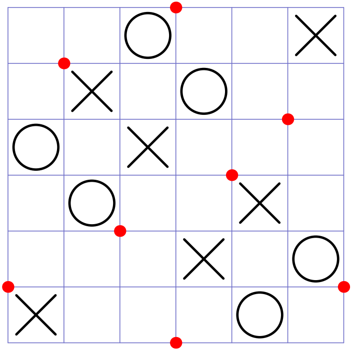

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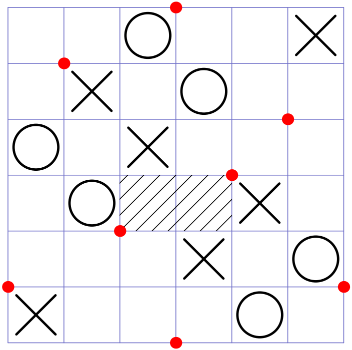

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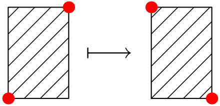

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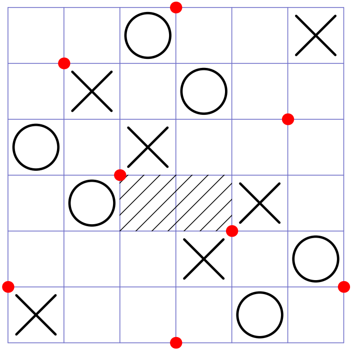

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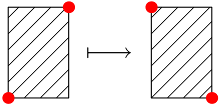

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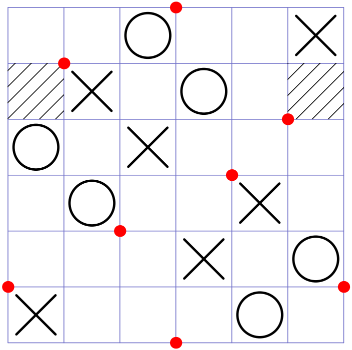

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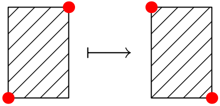

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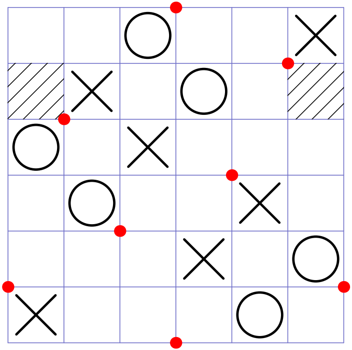

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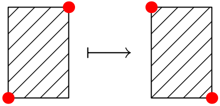

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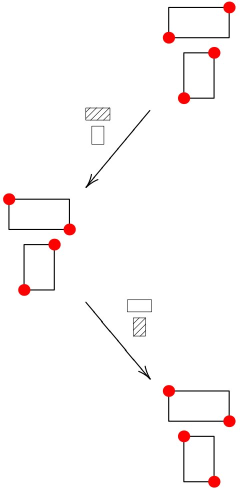

Each term in ∂^2 must have a mate:

- If rectangles are disjoint, take rectangles in either order.
- If rectangles share a corner, decompose the union in another way.

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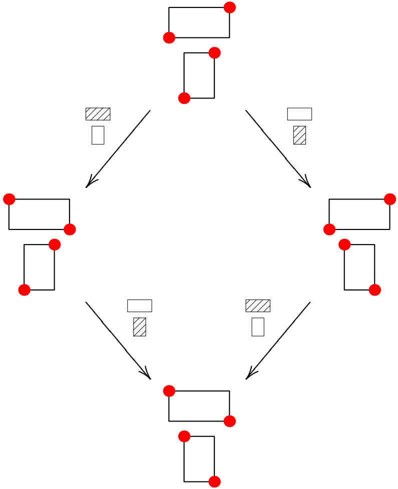

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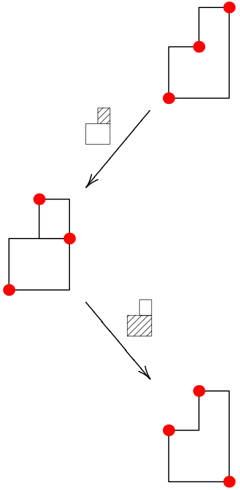

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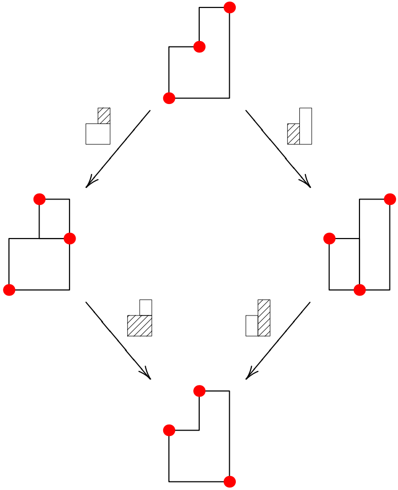

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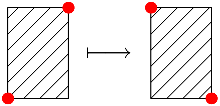
Computing *HFK*: Gradings on $\widetilde{CK}$

In the plane,



removes one *inversion*.

For $A, B, C \subset \mathbb{R}^2$,

$$\mathcal{I}(A, B) := \#\{\, {}_a\square^b \mid a \in A, b \in B \,\}$$
$$\mathcal{I}(A - B, C) := \mathcal{I}(A, C) - \mathcal{I}(B, C)$$

For $\mathbf{x}$ a generator, $\mathbb{X}$ the set of X's, $\mathbb{O}$ the set of of O's, the gradings are:

- **Maslov:** $M(\mathbf{x}) := \mathcal{I}(\mathbf{x} - \mathbb{O}, \mathbf{x} - \mathbb{O}) + 1.$
- **Alexander:**
  $A(\mathbf{x}) := \frac{1}{2}\big(\mathcal{I}(\mathbf{x} - \mathbb{O}, \mathbf{x} - \mathbb{O}) - \mathcal{I}(\mathbf{x} - \mathbb{X}, \mathbf{x} - \mathbb{X}) - (n-1)\big).$

Computing *HFK*: The answer

Theorem (Manolescu-Ozsváth-Sarkar)

For G a grid diagram for K,

$$H_*(\widetilde{CK}(G)) \simeq \widehat{HFK}(K) \otimes V^{\otimes n-1}$$

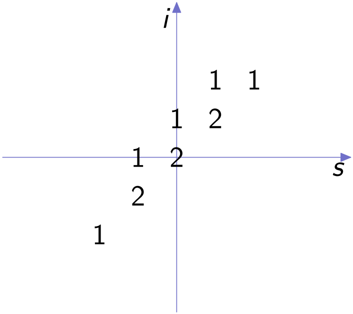
where $V := (\mathbb{Z}/2)_{0,0} \otimes (\mathbb{Z}/2)_{-1,-1}$.

Gillam and Baldwin used this to compute $\widehat{HFK}$ for all knots with ≤ 11 crossings, including new values of knot genus.

Outline

Improving the answer

$\dim \widehat{HFK}_i(K;s)$:



To remove factors of $V^{\otimes n-1}$:

HFK^-: variant of $\widehat{HFK}$
Module over $\mathbb{Z}/2[U]$
U has degree $(-1,-2)$
Related to $\widehat{HFK}$ by Univ. Coeff. Thm.
To compute: Add one U_i for each O

Complex $CK^-(G)$ over $\mathbb{Z}/2[U_1,\dots,U_n]$
∂ counts rects. that contain only O's, weighted by corresponding U_i.

Theorem (Manolescu-Ozsváth-Sarkar)

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Improving the answer

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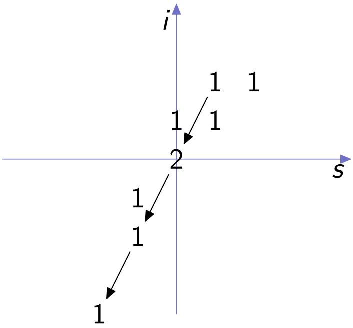

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Further variants

Can also:

- Allow rectangles to cross X's to get a filtered complex, and
- Add signs (in essentially unique way) to work over $\mathbb{Z}[U]$.

Outline

Combinatorial invariance

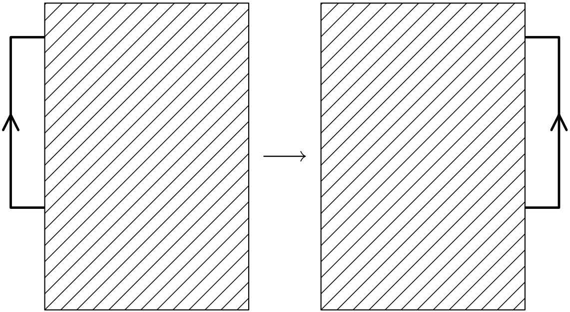
Theorem (Manolescu-Ozsváth-Szábo-T.)

For any sequence of elementary grid moves, there is an explicit chain map exhibiting invariance of HFK^-.

Conjecture (Naturality or Functoriality)

The chain map depends only on isotopy class of sequence of elementary grid moves. That is, oriented mapping class group of K acts on $HFK^-(K)$.

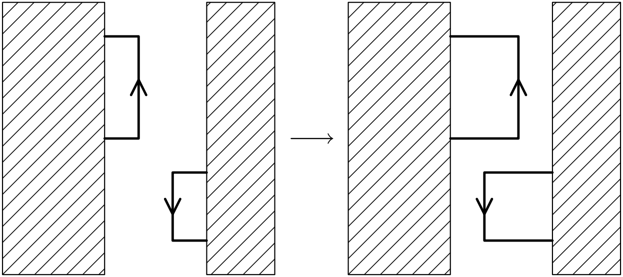
Elementary grid moves



- **Cycle:** Move left column to right, or top row to bottom.
- **Commute:** Switch two non-interfering columns or rows.
- **Stabilize:** Introduce a notch at a corner.

(Cromwell '95, Dynnikov '06)

Elementary grid moves



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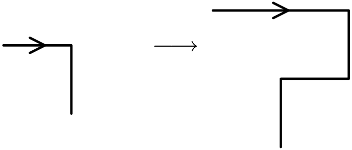

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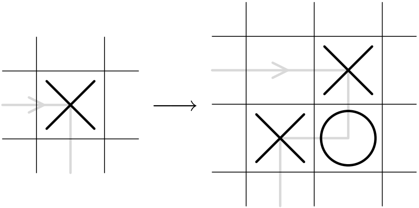

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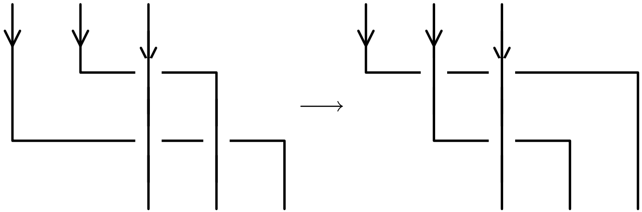

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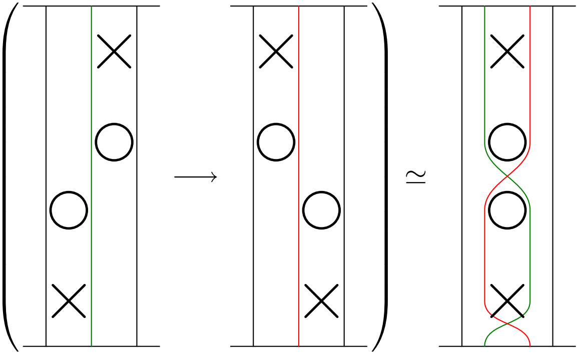
Where's Reidemeister III?

(Cromwell '95, Dynnikov '06)

Chain map for commutation counts pentagons



To construct a chain map for commutation, draw two versions of the middle gridcircle on a single diagram.

The chain map counts empty pentagons going between the two gridcircles.

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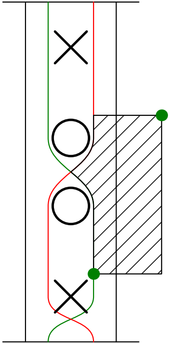

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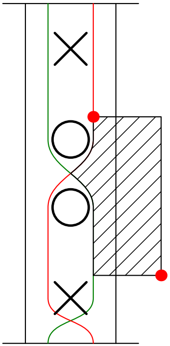


To construct a chain map for commutation, draw two versions of the middle gridcircle on a single diagram.

The chain map counts empty pentagons going between the two gridcircles.

Outline

Introduction

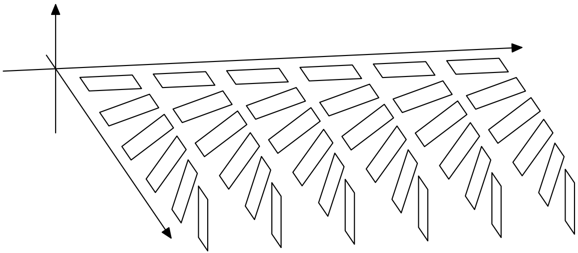
Computing *HFK*

Variants

Grid moves

- **Transverse knots**

Contact structures and knots



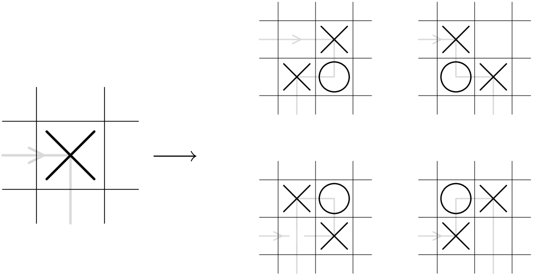
A *contact structure* is a twisted 2-plane field:
if α is a 1-form defining the plane field, $\alpha \wedge d\alpha$ is positive.
(Warning: above contact structure is reversed.)

A *Legendrian knot* is a knot that is tangent to the plane field.
A *transverse knot* is a knot that is transverse to the plane field.

Transverse knots have one easy invariant, the *self-linking number*.

Question. Can we find transverse knots with the same classical knot type and self-linking number?

Ways to stabilize

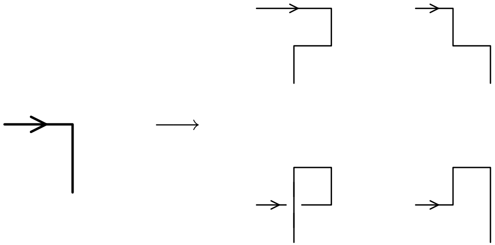


Four ways to stabilize: Where to leave the empty square?

- Two diagonal opposite ways preserve Legendrian knot.
- Two adjacent ways preserve closed braid.
- Three ways preserve transverse knot.

Warning: The Legendrian/transverse knots are mirrored.

Ways to stabilize



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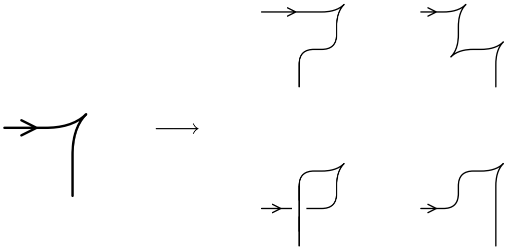

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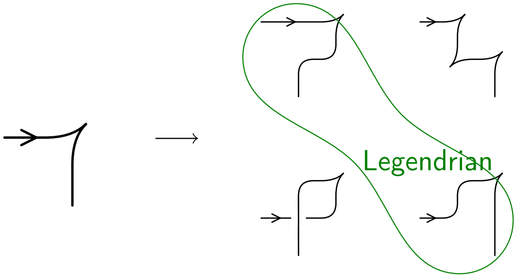

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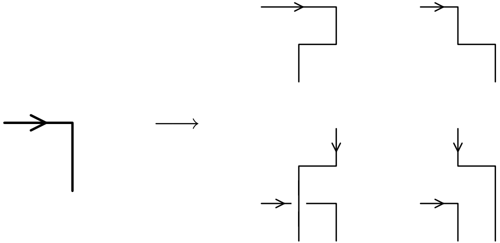

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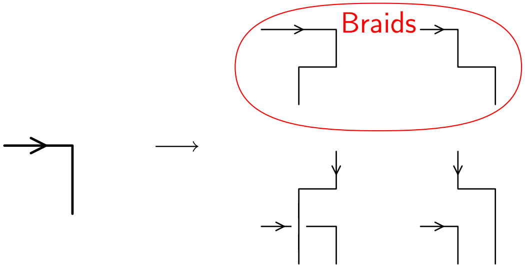

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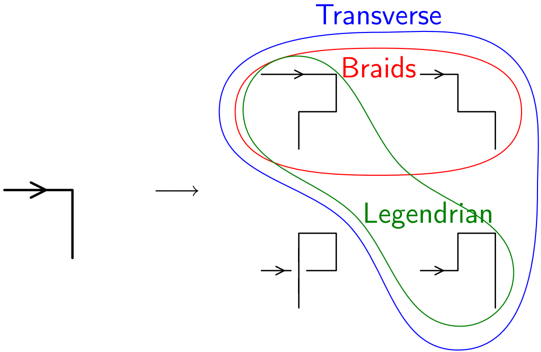

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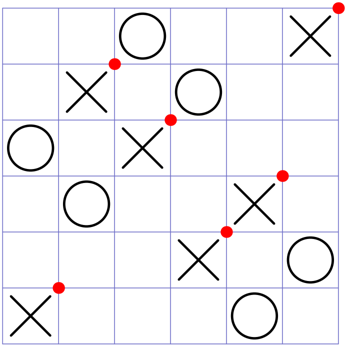
Transverse invariant: Definition



Definition

The *canonical generator* $\mathbf{x}^+(G)$ is given by the upper-right corner of each X.

Facts:

- $\partial \mathbf{x}^+ = 0$. (The X's block any rectangles.)
- $[\mathbf{x}^+(G)]$ maps to $[\mathbf{x}^+(G')]$ under commutation and 3 out of 4 stabilizations.

Theorem (Ozsváth-Szabó-T.)

$[\mathbf{x}^+(G)]$ in $HFK^-(m(K))$ is an invariant of the transverse knot represented by G, up to quasi-isomorphism of filtered complexes.

Transverse invariant: Properties

Let G be a grid diagram representing the transverse knot $\mathcal{T}$.

- $\mathbf{x}^+(G)$ lives in bigrading $(s, 2s)$, where $s = \frac{sl(\mathcal{T})+1}{2}$.
- If $\mathcal{T}'$ differs from $\mathcal{T}$ by a positive stabilization, then $[\mathbf{x}^+(\mathcal{T}')] = U[\mathbf{x}^+(\mathcal{T})]$.
- $[\mathbf{x}^+(\mathcal{T})] \neq 0$ in HFK^-.

Corollary

For any transverse knot $\mathcal{T}$ of topological type K,

$$\frac{sl(\mathcal{T})+1}{2} \leq \tau(K) \leq g_4(K)$$

where $\tau(K)$ is the largest Alexander grading which has an element which is not U torsion.

Transverse invariant: Examples

Let $\theta(\mathcal{T})$ (resp. $\widehat{\theta}(\mathcal{T})$) be the transverse invariant in $HFK^-(m(K))$ (resp. $\widehat{HFK}(m(K))$).
$\widehat{\theta}(\mathcal{T}) = 0$ iff $\theta(\mathcal{T})$ is divisible by U.

Theorem (Ng-Ozsváth-T.)

The knots $m(10_{132})$ and $m(12n_{200})$ have two trans. reps. with same sl, one with $\widehat{\theta} = 0$ and one with $\widehat{\theta} \neq 0$.

This technique also works for the $(2,3)$ cable of the $(2,3)$ torus knot, originally found by Etnyre-Honda and Menasco-Matsuda.

Let δ_1 be the next differential in the spectral sequence on $\widehat{HFK}$.

Theorem (Ng-Ozsváth-T.)

The pretzel knots $P(-4,-3,3)$ and $P(-6,-3,3)$ have two trans. reps. with same sl, one with $\delta_1 \circ \widehat{\theta} = 0$ and one with $\delta_1 \circ \widehat{\theta} \neq 0$.

Transverse invariant: Going further

Theorem (Ng-Ozsváth-T.)

If the Naturality Conjecture is true, then the twist knot 7_2 *has two trans. reps. with the same sl, with* $\widehat{\theta}$ *in different orbits of the mapping class group.*

But θ is not a complete invariant: Birman and Menasco have classified closed 3-braids up to transverse isotopy.
In their small examples of distinct transverse knots, θ lives in a 1-dimensional space, so cannot distinguish them.